RustFest Global 2020

Architect a High-performance SQL Query Engine in Rust

Jin Mingjian

Outline

- Introduction
- Related works
- Engineering in Rust
- System and Language
- Execution
- Evaluation
- Future insights

Introduction

- Jin Mingjian
- Ph.D., Data nerd
- Fields: bigdata engineering, high performance infrastructure, language and system
 - Cassandra, Spark, Impala, Kudu, ClickHouse...
 - jinmingjian.xyz/resume

Introduction

- For this talk
 - Works and practices in my recent open-source project TensorBase which is building a modern big data warehouse with Rust and its friend C
 - So, all working codes which covered or not covered in this talk are available in the site: tensorbase.io

Introduction

- For audiences of this talk
 - Keep content **as simple as possible**
 - Indeed both high performance and bigdata system are hard
 - Implementation details are ignored intentionally
 - All are available on the project sites
 - Ask questions, re-read this presentation or join the community

Related works

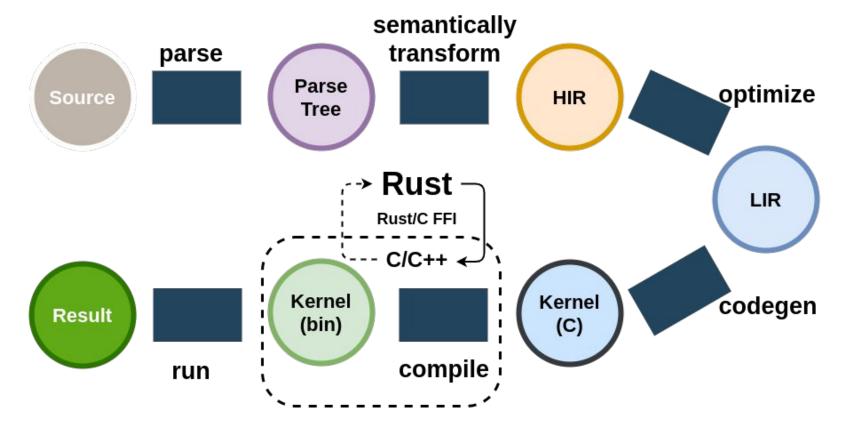
- Related projects in Rust
 - [DataFusion](<u>https://github.com/apache/arrow/tree/master/rust/datafus</u> ion)
 - Use (but not control) Apache Arrow
 - Far-from-performance-oriented architecture: Mixed Arrow and Rust operator impls (in AOT) (traditional design)
 - (opensource) Spark is a bottom baseline from the view of performance
 - The state-of-the-art speed is 100x faster than that of Spark

Related works

- Ruby Y. Tahboub et. al., "How to Architect a Query Compiler, Revisited"
 - Paper only but inspires TensorBase
 - Generative programming v.s. TensorBase's pipeline based IR and literal codegen
 - Generative representations on op-level are verbose and slow to compile
 - (generative) Style is **orthogonal to performance**

Related works

- Shoumik Palkar et. al., "Weld: A Common Runtime for High Performance Data Analytics"
 - Dedicated Weld IR for data vs TenosrBase's in-language IR
 - Abstract overhead
 - Deep binding to LLVM apis which is hard to upgrade without LLVM experts
- Problem of academic papers
 - Less engineering oriented
 - Hard to be validated by the public



- Engineering
 - **Engineering oriented language** design makes development agile and comfortable
 - TensorBase benefits from
 - Crates ecosystem, cargo, tooling/IDE, modules, macros, language(memory safe), language(zero-overhead C interop), language(ADT)
 - TensorBase: **one-person project** in several months (hope more to join in:)

- General aspects of TensorBase in engineering
 - Performance designed in core
 - Modernization
 - Bleeding edge nightly
 - Good practices
 - KISS
 - Minimization of Dependencies
 - Highly hackable

- Cargo
 - Highly extensible
 - Useful commands (except that most commonly used):
 - *add*: for add latest version of deps
 - **expand**: for proc macro debugging
 - **tree**: for transitive deps checking

- Proc macro
 - Problems
 - Learning curve is high
 - Too many out-of-date materials but less practices
 - Tooling is in primary stage
 - Hard to debug

- Proc macro (cont.)
 - Suggestions
 - use **nightly** as possible
 - enable feature nightly proc_macro_diagnostic, proc_macro_span for better debugging output
 - Proc_macro2 should improve the testability

- C interoperability
 - Zero overhead
 - Resource management
 - Objects in C is managed manually but not in safe Rust
 - Error handling
 - Too many "unsafe"s and "as"
 - Nice watch PR: [RFC Safer Transmute](/rfcs/pull/2981)

- Concurrency
 - Nice for share-nothing thread safety
 - Awkward when memory sharing needed
 - Memory sharing cornerstone of modern multicores
 - High frequency copy is a performance disaster
 - "Channel" is behind on sharing
 - Lack memory model like in Java/C++

- Concurrency (cont.)
 - E.g. Global (Safe) Sharing/Singleton

pub static CAT: Lazy<Mutex<Catalog>> = Lazy::new(
 let conf: Conf = Conf::load(load_path: None).
 let schema_dir: String = conf.schema.schema_d

Question: Oneshot change to the global but not at the declaration point?

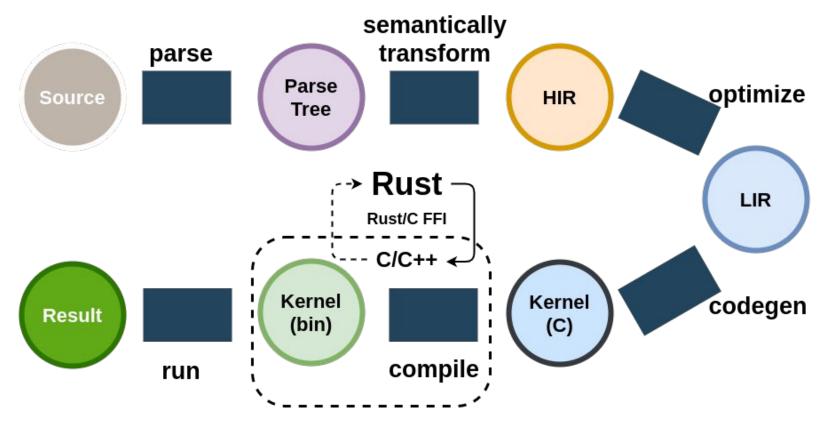
Safe Rust: Lazy lock

Comment: Locks are heavy. It could be safe if have change(write) "happens before" use(read) (common high-perf pattern in Java)

- Concurrency (cont.)
 - Async-await
 - Style is orthogonal to performance
 - Great concept but too deep boxing in kinds of implementations
 - Hard to debug when something wrong
 - Not used in TensorBase now

- Lifetime
 - Engineering excellence but make codes complex
 - Always recommends: to dance with, rather than to evade
 - Compiler enforces more than that needed when not smart
 - Alternatives
 - Arena allocator: TensorBase IR
 - Unsafe into C: TensorBase kernel algorithms in C







- Input
 - Query: Plain SQL
 - Not necessarily SQL
- Parse Tree
 - PEG, based on [Pest](https://github.com/pest-parser)
 - Lexical validation
 - (free-style) AST



- IR (Intermediate Representation)
 - Layered
 - No reinventing
 - Reuse modern low-level compilation infra as possible
 - Many optimizations in popular engines not needed any more
 - e.g. CSE(Common Subexpression Elimination)
 - select (a+123) as c1, (a+123)*2 as c2 from ...



- HIR (High-level IR)
 - Data related optimizations which can be not handled by low levels compilers
 - Semantic validation
 - Relational Algebra
 - e.g. predicate pushdown
 - Some RAs can be optimized by low-level compiler



- HIR (High-level IR) (cont.)
 - Unified RA operators
 - Core: 4 ops
 - ->(map), +(union/agg), *(join), <>(sort/top)
 - Inspired by J. Kepner's associative array
 - In Rust...



Prettyprinted Dump of HIR





- "Sea of pipes"
 - Unify data and control-flow dependencies in graph of "pipes"
- Pipe (a.k.a. Pipeline) (note: back to the dump)
 - Operator-fused computing/data unit (being extended to more)
 - Operator-level volcano model is systematically low inefficient
 - Unified dual view Data and Op
 - No more pull or push style, no more data center or control center



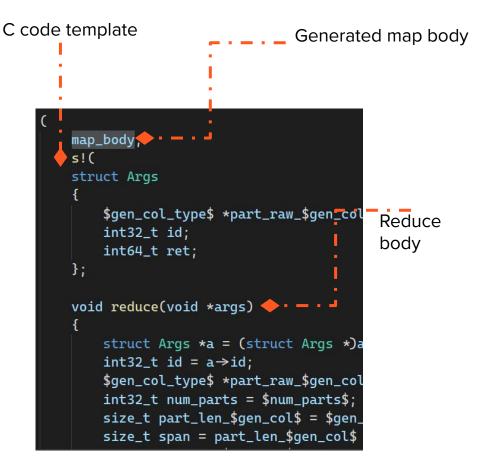
- LIR (Low-level IR)
 - Low level optimizations which can be not handled by low level compilers and not conveniently handled by high level logics
 - Platform related semantics
 - e.g. multi-cores + codegen
 - Boundaries between the high and low are not fixed
 - Scheduling will join later



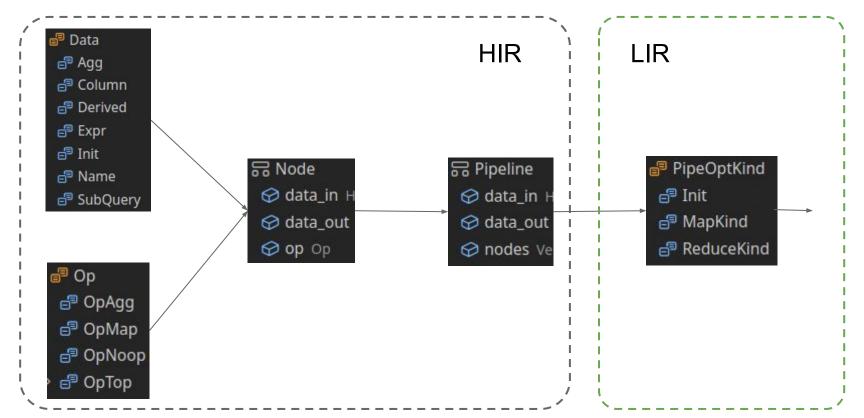
- LIR (Low-level IR) (cont.)
 - Parallelization representation for multi-cores
 - Generic DAG scheduling has abstract overhead
 - TensorBase: enable parallelism patterns
 - Map/reduce, fork-join...
 - Linearization representation for codegen
 - Proc macro template: provide human readable linear mapping from LIR to C primitives

Language

- LIR (Low-level IR) (cont.)
 - Right C template enables a simple map reduce pattern
 - Free style human oriented string interpolation
 - Runnable and debuggable in IDE with some setups







Execution

- Core
 - Decentralized, self-scheduling, JIT compilation based kernel
 - Decentralized != embarrassingly parallel
 - v.s. Popularly centralized scheduling
 - Why JIT compilation
 - Top performance
 - Arch agnostic

Execution

- Engine
 - Modified Clang based C JIT compiler
 - vs. popular LLVM IR based JIT engine(ClickHouse, Apache Impala...)
- Why C
 - Run on (almost) everything (CPU/GPU/FPGA/Accel)
 - Human debuggable
 - Fast enough compilation for OLAP

Execution

int64_t kernel()

```
char *blk_raw_c0 = NULL;
char fpath[64];
sprintf(fpath, "/data/n3/data/%d", 0); //TEMP
ker_scan(&blk_raw_c0, fpath);
```

```
struct Args s[48];
pthread_t ths[48];
for (size_t i = 0; i < 48; i++)
{
    s[i].id = i;
    s[i].part_raw_c0 = blk_raw_c0;
    pthread_create(&ths[i], NULL, reduce, &s[i]);
}
for (size_t i = 0; i < 48; i++)
{
</pre>
```

TensorBase generated kernel

// main function int main(int x0, char** x1) { long x2 = DEFAULT_INPUT_SIZE; long x3 = x2;long x4 = 0L;char** x7 = (char**)malloc(x2 * sizeof(char*)); **char**** x8 = x7: int* x9 = (int*)malloc(x2 * sizeof(int)); int* x10 = x9:long x14 = 0L;bool x15 = false: int x11 = open("Emp.tbl",0); long x12 = fsize(x11);char* $x13 = mmap(0, x12, PROT_READ,$ MAP_FILE | MAP_SHARED, x11, 0); bool x101 = !true; for (;;) { // ... parse and load data elided ... **char**** x96 = x8: char* x53 = x13+x42; x96[x76] = x53;

V.S.

Ruby Y. Tahboub et. al., Sigmod18

** TensorBase M0 benchmark: https://tensorbase.io/

Phase	Language (TensorBase M0)	Time	Ref Time (from public)	Language (Ref system)
Parsing (TPC-DS)	Rust	130 us**	~50 us*	C++
Parsing/IR/codegen (one column sum)	Rust	130 us**	~1 ms*	C++
C Kernel JIT Compilation	C/C++	13 ms (boost) - 20 ms (no boost) (Q#1 like, -O2)**	59 ms* (TPC-H Q#1, opt.)	C++
End-to-end Query Time (one column sum with 1.47B row NYC taxi dataset)	Almost in Rust (mixed)	~ 60 ms (compilation cached)- ~100 ms(compilation uncached)**	642 ms** In ClickHouse 20.05 (compilation cached)	C++

- Points (part 1)
 - Rust is lighting fast even untuned (but not Rust compilation)
 - C based JIT Compilation is lighting fast even untuned (on par with LLVM IR)
 - C based JIT Compilation is much faster than C++ (and Rust) even untuned
 - C based JIT Compilation is quite enough for OLAP even untuned

- Points (part 2)
 - Saturates memory bandwidth in core of such runs
 - Can't not be faster in single 6-channel xeon sp with 100GB/s memory bandwidth(measured by vtune) for memory bound applications
 - Napkin math: 1.47*4/100 = ~0.06 sec (= 60msec)
 - Simple sum aggregation query
 Tight loop in kernel and finally
 vectorized by compiler

```
size_t s = 0;
for (size_t i = 0; i < blk_len_c0; i++)
{
    int32_t c0 = blk_c0[i];
    s += c0;
} You, 3 months ago • base m0</pre>
```

- Points (part 3)
 - Partial compilation
 - Makes compilation time of TensorBase is correlated to the size of hot kernel (rather than the total size of execution codes)
 - Lower bound of LLVM compilation overhead is high
 - -O0: ~9ms for helloworld (untuned)
 - Too many passes
 - Speed does not come for free

C	23.1%)	X86 DAG->DAG Instruction Selection
C	7.8%)	Loop Strength Reduction
C	7.1%)	Induction Variable Simplification
(5.2%)	Loop Vectorization
C	4.2%)	Machine Instruction Scheduler
C	3.8%)	SROA
C	3.6%)	Unroll loops #2
(3.5%)	Combine redundant instructions
C	3.3%)	Combine redundant instructions #8
C	2.0%)	Simplify the CFG
C	1.7%)	Induction Variable Users
C	1.5%)	Combine redundant instructions #6
C	1.0%)	
C	0.9%)	Loop Load Elimination
C	0.9%)	Combine redundant instructions #3
C	0.8%)	
C	0.8%)	
C	0.8%)	
(0.8%)	SLP Vectorizer

- Storage layer
 - Popular storage and compute separation is genetically less efficient
- Optimizer
 - Makes the queries which can not be optimized fastest
 - Data driven, low entropy inference (e.g. category theory)
- Execution engine

- Tiered C compilation
 - OLTP style wants faster codegen
 - C compilation/interpretation possibly done in microseconds or less
- Alternative JIT compilation choices
 - e.g. Rust + Cranelift codegen (JIT) backend
- Scheduling

- Modern hardware oblivious
- Distributed
 - Techniques on single node does *NOT* mean they only works for single node
 - Consistency as plugin
- Engineering
 - Contracts and formal verifications
 - Rust C Rust chaining and maximized engineering Rust

- TensorBase 2020.11 (WIP)
 - Main operators on on single table
 - Storage Layer v1
 - Compatibility and hybrid deployment with ClickHouse
 - Compatible to ClickHouse Native Protocol
 - Partially compatible to ClickHouse on-disk data storage (read-only)

- TensorBase 2020.11 (WIP) (cont.)
 - Superb in complex aggregations (e.g. group by)
 - Early results (compared to ClickHouse)
 - 6x faster in several "group by" queries on top of ClickHouse MergeTree storage with a real-world billion level dataset
 - 10x faster if do some simple optimizations
 - Faster results expected when TensorBase own storage coming



- Abstract overhead everywhere
 - Carefully make trade-offs
- High-performance programming paradigm in Rust
 - Safe(almost) + unsafe (bounded but more freedom)
- **Top performance** OLAP is firstly achieved with **engineering Rust**
 - All shown can be picked up from the open source project (tensorbase.io)

